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(54) Abstract Title
Apparatus for generating a session key to encrypt or decrypt messages

(57) An apparatus for encrypting and decrypting messages is arranged to randomly generate a session key of a variable selected number of characters. The randomly generated characters are distributed in sequence into a predetermined number of groups (in the same manner as dealing a pack of cards out) to form a set of primitives. Successive pairs of primitives are combined in a XOR procedure and a successive multiple of 100 is added to each result. The primitives thus produced are used in accordance with a predetermined algorithm to form a cypher key stream for encrypting or decrypting successive characters of a message. The apparatus is for use in a facsimile machine.

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At least one drawing originally filed was informal and the print reproduced here is taken from a later filed formal copy.

The claims were filed later than the filing date but within the period prescribed by Rule 25(1) of the Patents Rules 1995.

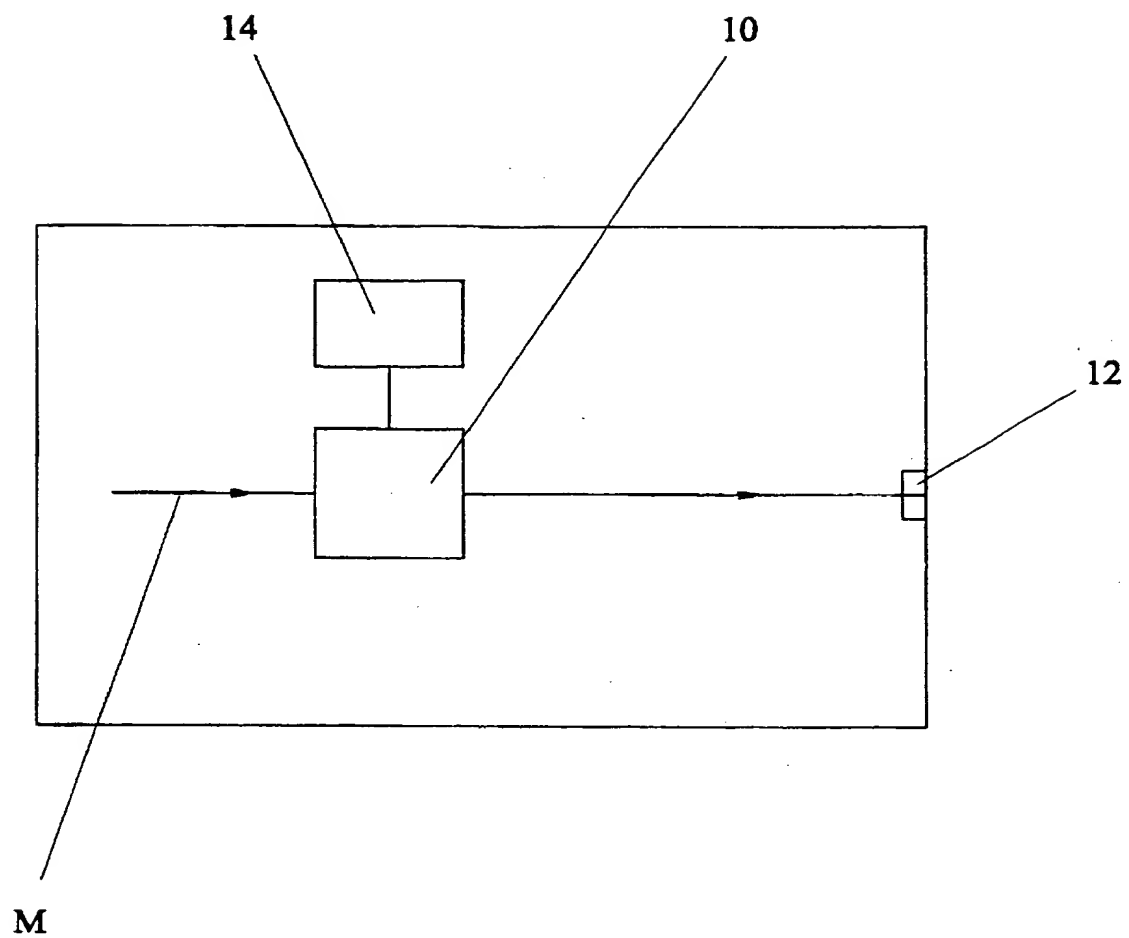


FIG. 1

- (1) Randomly generated 56 figures secret key
 44490925319354895003213478116711124821791736922366044359 = 186 bits
- | | | | | | | | | | | | | | |
|------|------|------|-------|-------|------|------|------|------|------|------|-------|-------|-------|
| 4863 | 4976 | 4519 | 9012 | 0012 | 9323 | 2246 | 5186 | 3320 | 1414 | 9774 | 3893 | 5515 | 4179 |
| 4863 | 399 | 4136 | 13084 | 13072 | 6011 | 8125 | 3071 | 1799 | 641 | 9391 | 11162 | 14433 | 10290 |
| 4963 | 599 | 4436 | 13848 | 13572 | 6611 | 8825 | 3171 | 1999 | 941 | 9791 | 11662 | 15033 | 10990 |
- (2) Randomly generated 48 figure secret key
 019301259141828933881545855137151100195663513916 = 159 bits
- | | | | | | | | | | | | | | |
|------|------|------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| 0835 | 1971 | 9313 | 3359 | 0811 | 1816 | 210 | 550 | 941 | 159 | 485 | 156 | 856 | 213 |
| 835 | 1264 | 8337 | 11662 | 11941 | 10685 | 10607 | 11081 | 10468 | 10363 | 10654 | 10498 | 10842 | 10895 |
| 935 | 1464 | 8637 | 12062 | 12441 | 11285 | 11307 | 11181 | 10668 | 10663 | 11054 | 10998 | 11442 | 11595 |
- (3) Randomly generated 40 figure secret key
 8272611489892272004839562205012958050601 = 133 bits
- | | | | | | | | | | | | | | |
|-----|------|-----|-----|------|------|------|-----|-----|-----|-----|------|------|------|
| 870 | 221 | 702 | 209 | 645 | 188 | 130 | 495 | 850 | 966 | 820 | 921 | 20 | 25 |
| 870 | 955 | 261 | 468 | 849 | 1005 | 879 | 640 | 466 | 532 | 288 | 697 | 685 | 692 |
| 970 | 1155 | 561 | 868 | 1349 | 1605 | 1579 | 740 | 666 | 832 | 688 | 1197 | 1285 | 1392 |
- (4) Randomly generated 32 figure secret key
 71936502524388145097469032024067 = 106 bits
- | | | | | | | | | | | | | | |
|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| 714 | 140 | 956 | 307 | 69 | 57 | 04 | 26 | 59 | 20 | 43 | 32 | 80 | 82 |
| 714 | 582 | 506 | 201 | 140 | 181 | 177 | 171 | 144 | 132 | 175 | 143 | 223 | 141 |
| 814 | 782 | 806 | 601 | 640 | 781 | 877 | 271 | 344 | 432 | 575 | 643 | 823 | 841 |
- (5) Randomly generated 18 figure secret key
 241895746918348060 = 60 bits
- | | | | | | | | | | | | | | |
|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| 20 | 40 | 16 | 80 | 9 | 5 | 7 | 4 | 6 | 9 | 1 | 8 | 3 | 4 |
| 28 | 52 | 36 | 116 | 125 | 120 | 127 | 123 | 125 | 116 | 117 | 125 | 126 | 122 |
| 128 | 252 | 336 | 516 | 625 | 720 | 827 | 223 | 325 | 416 | 517 | 625 | 726 | 822 |
- (6) Randomly generated 12 figure secret key
 449078277840 = 40 bits
- | | | | | | | | | | | | | | |
|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| 4 | 4 | 9 | 0 | 7 | 8 | 2 | 7 | 7 | 8 | 4 | 0 | 4 | 4 |
| 4 | 0 | 9 | 9 | 14 | 6 | 4 | 3 | 4 | 12 | 8 | 8 | 12 | 8 |
| 104 | 200 | 309 | 409 | 514 | 606 | 704 | 103 | 204 | 312 | 408 | 508 | 612 | 708 |

FIG. 2

ENCRYPTION AND DECRYPTION KEY ARRANGEMENTS

The present invention relates to apparatus arranged to encrypt messages or decrypt messages, particularly to communications apparatus arranged to encrypt messages prior to transmission and decrypt received messages.

5 It is known to provide communications apparatus (for example facsimile machines) with the ability to encrypt messages prior to transmission and decrypt received messages. However, each such apparatus operates with a cypher of a predetermined, fixed cryptographic strength: two apparatus can
10 only communicate with each other if they both use cyphers of the same strength. There are many circumstances in which this limits the ability for communications to be established.

 We have now devised an arrangement in which the cypher can be varied in strength, so that when any two apparatus wish
15 to communicate with each other, a common cypher strength can be selected.

 Thus, in accordance with the present invention, there is provided an apparatus which is arranged to encrypt or decrypt messages, the apparatus being arranged to generate a
20 session key of a variable selected number of characters and to distribute the characters of said session key in sequence into a predetermined number of groups to form a corresponding predetermined number of primitives, and further arranged to use said primitives, in accordance with a predetermined algorithm,
25 to form a cypher key stream the characters of which are used in sequence to encrypt or decrypt successive characters (or other elements) of a message.

 In use of this apparatus, the length (i.e. the number of characters) of the session key can be selected: the longer
30 the session key, the greater will be the strength of the cypher.

 Once the length of the session key to be used is decided upon, the session key is preferably randomly generated.

 Preferably the characters (typically numerical
35 characters) of the session key are distributed into the predetermined number of groups in a manner forming a

corresponding set of multi-digit numbers. For example, the first term allocated to each group may form the first digit of a multi-digit number, the second term allocated to that group forms the second digit of the multi-digit number, and so on.

5 Preferably these multi-digit numbers are processed further in order to produce the corresponding set of primitives, used to form the cypher key stream.

 Preferably successive pairs of these multi-digit numbers are then subjected to an XOR (exclusive OR) process to
10 form a corresponding set of results.

 Preferably predetermined values are then added to the respective results of the XOR process, to form a corresponding set of primitives. Preferably different values are added to the different results of the XOR process: preferably these
15 different values are different multiples of a basic value. For example, 100 may be added to the first XOR result, 200 to the second, and so on.

 An embodiment of the present invention will now be described with reference to the accompanying drawings, in
20 which:

 FIGURE 1 is a schematic block diagram showing part of the electronic system of communications apparatus in accordance with the present invention; and

 FIGURE 2 is a table showing the formation of six
25 different groups of primitives from six session keys of different lengths.

 Referring to Figure 1, a communications apparatus (e.g. a facsimile machine) comprises means 10, in the form of a microprocessor, for encrypting a plain message M prior to
30 transmission via a port 12. The microprocessor 10 is provided with a program memory 14 which stores an encryption algorithm and also an algorithm for forming a group of primitives from a session key. The microprocessor is able to generate a session key on a random basis, of selected length. The
35 microprocessor is also arranged to correspondingly decrypt messages received via the port 12.

 In effecting communication between two apparatus, these follow an initial protocol to determine the cryptographic strength to be employed: this determines the length of the

session key to be used. Then the session key is randomly generated by the microprocessor 10 in one of the apparatus: Figure 2 shows six different examples, in which session keys of 56, 48, 40, 32, 18 and 12 decimal digits (186, 159, 133, 106, 60 and 5 40 binary bits) are generated.

Once the session key of selected length has been generated, the microprocessor distributes its digits, one-after-another, into 14 groups, in the same manner as dealing a pack of cards out to the players of a card game. Thus, 10 referring to the first example in Figure 2, the first 14 digits (44490925319354) form the first digits of respective 4-digit numbers: continuing, the next 14 digits of the session key (89500321347811) form the second digits of the respective 4-digit numbers, the next 14 digits of the session key 15 (67111248217917) form the third digits of the respective 4-digit numbers and the final 14 digits of the session key (36922366044359) form the fourth (and final) digits of the respective 4-digit numbers. In the first example in Figure 2, 14 groups of 4-digit numbers are thus formed: however, in each 20 of the other examples, the number of digits in the session key is not divisible by the number of groups (14), so that 14 numbers of differing numbers of digits are formed (in some cases, only a single digit).

In the next step, the microprocessor 10 combines 25 successive pairs of the 14 numbers in an XOR (exclusive OR) procedure: in each of the examples shown in Figure 2, the second line gives the corresponding results. In particular, each number in the first line is combined with the XOR result of the proceeding number, in a process which involves an XOR 30 function or their binary equivalents.

In the next step (third line of each example shown in Figure 2), the microprocessor 10 adds a multiple of 100 to each of the 14 results formed by the XOR procedure. Thus, to the first result, 100 is added: to the second result, 200 is 35 added; to the third result, 300 is added, and so on up to the seventh result, to which 700 is added. Then, to the eighth result, 100 is added: to the ninth result, 200 is added, and so on up to the fourteenth result, to which 700 is added. The final results (last line in each of the 6 examples set out in

Figure 2) provide a set of 14 primitives.

It will be appreciated that the second and third steps which have been described add complexity to the primitives finally produced. The third step in particular ensures that
5 none of the primitives will be zero.

The 14 primitives thus produced are used by the microprocessor, in accordance with the encryption algorithm, to form a cypher key stream comprising a long stream of digits. Then, in order to encrypt a plain message, the digits of this
10 stream are taken one-after-another, and used in accordance with an encryption algorithm to encrypt respective, successive elements (e.g. characters or groups of characters) of the message to be transmitted. Similarly, in order to decrypt a received message, the digits of the cypher key stream are taken
15 one-after-another and used, in accordance with a decryption algorithm (being the inverse of the encryption algorithm) to decrypt respective, successive elements of the received message.

Claims

- 1) An apparatus which is arranged to encrypt or decrypt messages, the apparatus being arranged to generate a session key of a variable selected number of characters and to
5 distribute the characters of said session key in sequence into a predetermined number of groups to form a corresponding predetermined number of primitives, and further arranged to use said primitives, in accordance with a predetermined algorithm, to form a cypher key stream the characters of which are used
10 in sequence to encrypt or decrypt successive characters (or other elements) of a message.
- 2) An apparatus as claimed in claim 1, arranged to generate said session key in random manner.
- 3) An apparatus as claimed in claim 1 or 2, arranged so
15 that the characters of the session key are distributed into said predetermined number of groups in a manner forming a corresponding set of multi-digit numbers.
- 4) An apparatus as claimed in claim 3, arranged to further process said multi-digit numbers to produce said primitives.
- 20 5) An apparatus as claimed in claim 4, arranged to subject successive pairs of said multi-digit numbers to an exclusive OR process to form a corresponding set of results, and to process said results to produce said primitives.
- 6) An apparatus as claimed in claim 5, arranged so to add
25 predetermined values to the respective said results, to form said primitives.
- 7) An apparatus as claimed in claim 6, arranged to add different said values to different said results.
- 8) An apparatus as claimed in claim 7, arranged such that
30 said different values are different multiples of a basic value.

9) An apparatus which is arranged to encrypt or decrypt messages, the apparatus being substantially as herein described with reference to the accompanying drawings.



Application No: GB 9814003.1
Claims searched: All

Examiner: Gareth Griffiths
Date of search: 13 August 1999

Patents Act 1977
Search Report under Section 17

Databases searched:

UK Patent Office collections, including GB, EP, WO & US patent specifications, in:

UK Cl (Ed.Q): H4P (PDCSP)

Int Cl (Ed.6): H04L 9/12, 9/18, 9/22, 9/26

Other: Online Databases: WPI, EPODOC, JAPIO

Documents considered to be relevant:

| Category | Identity of document and relevant passage | Relevant to claims |
|----------|---|--------------------|
| A | GB2301266 A (WILLIAM YIN SHAW) | |

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| X | Document indicating lack of novelty or inventive step | A | Document indicating technological background and/or state of the art. |
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